ECS 20 – Fall 2021 – P. Rogaway

https://www.cs.wm.edu/~tadavis/cs243/ch05s.pdf

Number Theory

- 1. constant symbol: 0
- 2. predicate symbol: <
- 3. function symbol: *S* (1-ary) (successor function)
 - + (2-ary)
 - (2-ary)
 - *E* (2-ary) omit because it's not part of PA?

Always add equality (=), which is reflexive, symmetric, transitive.

Already rich enough to make powerful statements in number theory. Eg: **Fermat's Last Theorem**: no three positive integers *a*, *b*, and *c* satisfy the equation $a^n + b^n = c^n$ for any integer value of *n* greater than 2:

 $(\forall a) (\forall b) (\forall c) (\forall n) (a \ E \ n + b \ E \ n = c \ E \ n \to n > S(S(0)))$ (Show how to define > using <, =, and negation). Or, similarly, write **Goldbach's Conjecture** in this language of number theory: that every even number more than 2 is the sum of two primes.

PA axioms from where? Now finding <u>http://www.cs.toronto.edu/~sacook/csc438h/notes/page96.pdf</u>, which doesn't include 7, 8...

Axioms of arithmetic ("Peano arithmetic")(Giuseppe Peano, 1889)

- 1. $(\forall x) (S(x) \neq 0)$
- 2. $(\forall x)(\forall y)(S(x) = S(y) \rightarrow x = y)$
- 3. $(\forall x) (x + 0 = x)$
- 4. $(\forall x)(\forall y)(x + S(y) = S(x+y))$
- 5. $(\forall x) (x \cdot 0 = 0)$
- 6. $(\forall x)(\forall y)(x \cdot S(y) = x \cdot y + x)$
- 7. $(\forall x)(\forall y))(\forall c) (x < y \rightarrow x + c \le y + c)$
- 8. $(\forall x)(\forall y))(\forall c) (x < y \rightarrow x \cdot c \le y \cdot c)$
- 9. For all predicates P

 $(P(0) \land (\forall n)(P(n) \rightarrow P(n+1)) \rightarrow \land (\forall n)(P(n))$

Numbers and Induction

Not a 1st order property

Alternatively: If a **set** contains zero and the successor of every **<u>number</u>** is in the set, then the set contains the natural numbers. This form does not seem as directly useful.

Principle of mathematical induction Different statement

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To prove a proposition P(n) for all integers n \ge n_0:
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1) Prove P(n_0) (Basis)
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2) Prove that $P(n) \rightarrow P(n+1)$ for all $n > n_0$ (Inductive step) (Inductive assumption)

The above sounds slightly more general (because I let you start at n_0), but easily seen to be equivalent.

Also equivalent: "**strong**" form of induction: To prove a proposition P(n) for all integers $n \ge n_0$: 1) Prove $P(n_0)$ (Basis) 2) Prove that $(P(1) \land ... \land P(n)) \rightarrow P(n+1)$ for all $n > n_0$ (inductive step)

Again equivalent. Sometimes easier to apply. The stronger inductive assumption may make it easier to get the conclusion.

EXAMPLE 0: Prove that the sum of the first *n* integers n(n+1)/2. Do this in two ways, either: (a) pictorially; (b) by induction. But where does the formula come from? Options: (i) write a table; (ii) use a definite integral to help make a guess; (iii) solve a system of equations using values from table.

EXAMPLE 1: Prove that the sum of the **odd** integers 2 ... 2n-1 is n^2 1 + 3 + ... + (2n-1) = n^2 .

Basis: *n*=1, check

Inductive step:

$$1 + 3 + \dots (2n - 3) = (n - 1)^{2} + 2n - 1 = + 2n - 1 = n^{2} - 2n + 1 + 2n - 1 = 1 = n^{2}$$

EXAMPLE 2: Use induction to prove that n^2+n is always even (divisible by 2).

Basis: n=0: fine. Inductive step: Assume that the statement is true for n=k. Thus, k^2+k is even. That is, $k^2+k = 2j$ for some integer j. Now what about $(k+1)^2+(k+1) -$ is it necessarily even?? Expanding out, this expression is $k^2 + 2k + 1 + k + 1 = k^2 + 3k + 2 = k^2 + k + 2k + 2$ = 2j + 2k + 2 = 2(j+k+1)

so it is even

Now, write k^2+k^2 + k^2+k as part of an equation which denotes that it is divisible by 222.

EXAMPLE 3. Sam's Dept. Store sells envelopes in packages of 5 and 12. Prove that, for any $n \ge 44$, the store can sell you exactly n envelopes. [GP, p.147]

Basis: 44 = 2(12) + 4(5) 45 = 9(5) 46 = 3(12) + 2(5)2...?

SUPPOSE: It is possible to buy *n* envelopes for some $n \ge 44$. **SHOW**: It is possible to buy *n*+1 envelopes

x	х хх	xx :	к х	$\mathbf{x}\mathbf{x}$	x	\mathbf{x}	x	xx	хэ	xxxxx	xxxxx	xxxxxx
12345678	90123	456789	012	3450	6789	901	L23	45678	9012	3456789	901234	567890
0	1	:	2			3			4		5	6

• If purchasing at least 7 packets of 5: trade in seven packets of five for three packets of 12:

 $\begin{array}{c} 7(5) \rightarrow 3(12) \\ 35 \qquad 36 \end{array}$

If purchasing fewer than 7 packets of 5: i.e., purchasing at most 6 packets of 5, so at most 30 of the envelopes are in packets of 5; so what remains are ≥ 44 - 30 = 14 envelopes being bought in packets of 12, so ≥ 2 packets of twelve. So take 2 of the packets of 12 (i.e., 24 envelopes) and trade them for 5 packets of 5: 2(12) → 5(5) 24 25

EXAMPLE 4: Show that you can tile **any** "punctured" $2^n \times 2^n$ grid by *triominoes* <u>https://undergroundmathematics.org/divisibility-and-induction/triominoes/solution</u>

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(may be rotated)

Illustrate and prove, dividing board in into four $2^n \times 2^n$ to prove. Puncture the $2^{n+1} \times 2^{n+1}$ grid; tile that one of the four subgrids (by inductive assumption); puncturing three the three near-center center points (for the three $2^n \times 2^n$ pieces that lacking the puncture); recurse on those three pieces; add one more tromino.

EXAMPLE 5: Fundamental theorem of arithmetic. – Almost verbatim from <u>https://www.cs.wm.edu/~tadavis/cs243/ch05s.pdf</u> An example of "strong" induction".

Claim: if *n* is an integer greater than 1, then *n* can be written as the [unique] product of primes. Solution: Let P(n) be the proposition that *n* can be written as a product of primes. // Leave uniqueness for later, or omit it.

Basis: P(2) is true since 2 itself is prime.

Inductive step: The inductive hypothesis is P(j) is true for all integers j with $2 \le j \le k$. To show that P(k + 1) must be true under this assumption. Two cases need to be considered:

• If k + 1 is prime, then P(k + 1) is true.

• Otherwise, k + 1 is composite and can be written as the product of two positive integers *a* and *b* with $2 \le a \le b < k + 1$. By the inductive hypothesis *a* and *b* can be written as the product of primes and therefore k + 1 can also be written as the product of those primes. Hence, it has been shown that every integer greater than 1 can be written as the product of primes.

Uniqueness:

Let *N* be the *smallest* number that can be written in *two different ways* as the product of primes. Those two ways can have no prime in common or else we'd divide by it and have a smaller number that could be written in two different ways as the product of primes. Thus

$$N = p_1 p_2 \dots p_m = q_1 q_2 \dots q_n$$
$$= p_1 P = q_1 Q$$

Where $p_1 < q_1$ and the primes on the left, listed in increasing order, are disjoint from those on the right, also listed in increasing order. Now

$$(q_1 - p_1) Q < N$$

I claim that $(q_1 - p_1) Q$ is a multiple of p_1 , namely

because

$$p_1 P - p_1 O = N - p_1 O$$

 $p_1(P - Q) = (q_1 - p_1)Q$

and

$$(q_1 - p_1) Q = q_1 Q - p_1 Q = N - p_1 Q$$

By the unique factorization of number less than N we know that p_1 must occur in the factorization of $(q_1 - p_1)$ or in the factorization of Q.

- The first is impossible because if p_1 divides $q_1 p_1$ then it divides q_1 , but p_1 and q_1 are distinct primes.
- The second is impossible because the factors of *Q* were bigger than (as well as distinct from) *p*₁

Done.

The fundamental theorem of arithmetic is useful. We routinely like to think of numbers in terms of their prime factorization. Many questions become easier if presented a number in this form. Example:

How many factors does 360 have? First, write 360 in its prime factorized form: $360 = 2^3 \cdot 3^2 \cdot 5$ A factor must be of the form 2a 3b 5c where $a \in [0..3]$, $b \in [0..2]$, $c \in [0..1]$. So the number of factors 360 has is $4 \cdot 3 \cdot 2 = 24$.

If I asked you how many factors $2450250000 = 2^4 3^4 5^6 11^2$ has, you would answer $5 \cdot 5 \cdot 7 \cdot 3 = 525$

Is 1873215592 a square? No, because you can divide it by 2 three times, but no no more, so the factorization is $2^3 \cdot M$ where the prime factorization of *M* has no 2s in it; and a number is going to be a square iff all the powers of primes in the prime factorization are *even*.

EXAMPLE 6: Cake cutting

See <u>http://www.cs.berkeley.edu/~daw/teaching/cs70-s08/notes/n8.pdf</u> for a nice writeup

- 1. If *n* = 2, use the cut-and-choose protocol. Otherwise:
- 2. The first *n*–1 participants divide the cake by recursively invoking this procedure.
- 3. For i = 1,2,...,*n*−1, do:
 - a) Participant *i* divides her share into *n* pieces she considers of equal worth (by her measure).
 - b) Participant n collects whichever of those n pieces he considers to be worth most (by his measure).

Number of cuts:

 $T(n) = T(n-1) + (n-1)^2$

$$T(n) = T(n-1) + (n-1)^{2}$$

= T(n-2) + (n-1)^{2} + (n-2)^{2}
= T(n-3) + (n-1)^{2} + (n-2)^{2} + (n-3)^{2}
= T(n-3) + (n-1)^{2} + (n-2)^{2} + (n-3)^{2}
= 1 + 2^{2} + 3^{2} + 4^{2} + ... + (n-1)^{2}
approx. Integeral_1^n x2 approx n³/3

$$\sum_{i=1}^n i^2 \ = \ 1^2 + 2^2 + \dots + n^2 \ = \ rac{n(n+1)(2n+1)}{6}.$$

Prove by induction.

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$$(P(0) \land (\forall n)(P(n) \rightarrow P(n+1)) \rightarrow \land (\forall n)(P(n))$$

